

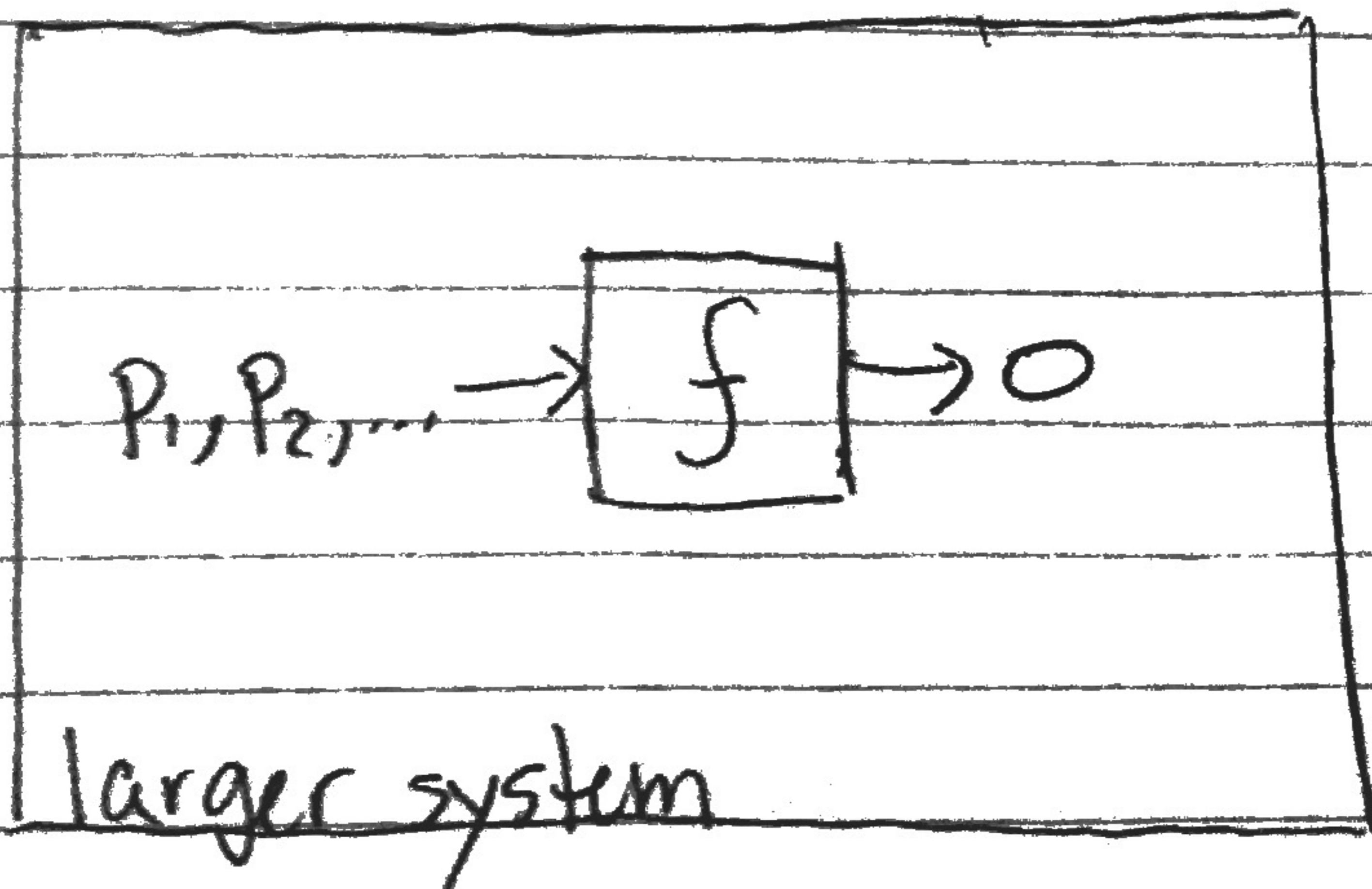
What is a pure function?

defn: A function f with input type A and output type B ($A \Rightarrow B$) is a computation that relates every value of $a \in A$ to exactly one value $b \in B$ such that b is determined solely by the value of a .

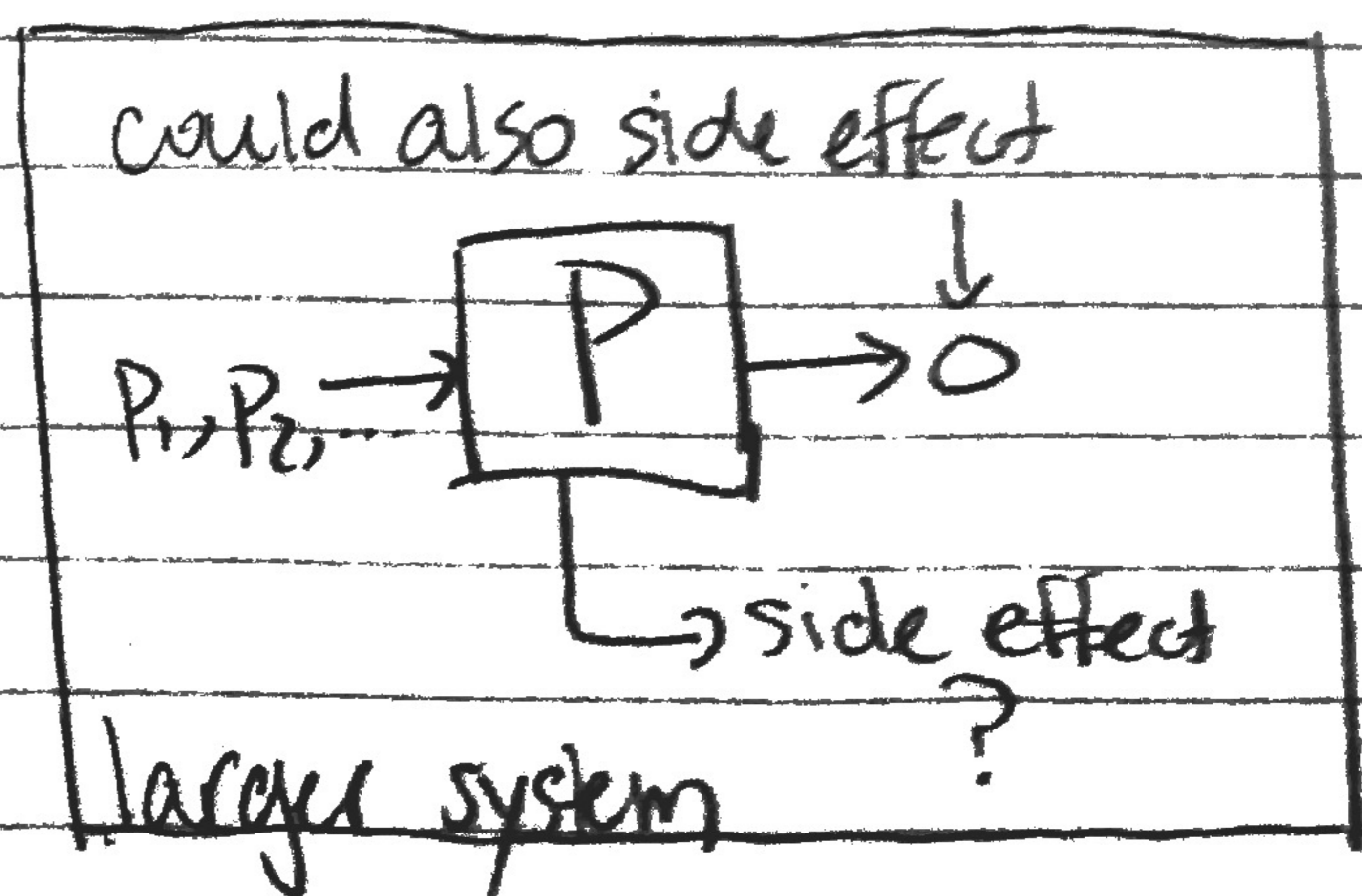


A function has no observable effect on the execution of a program other than to compute a result given its inputs.

(Pure) Function

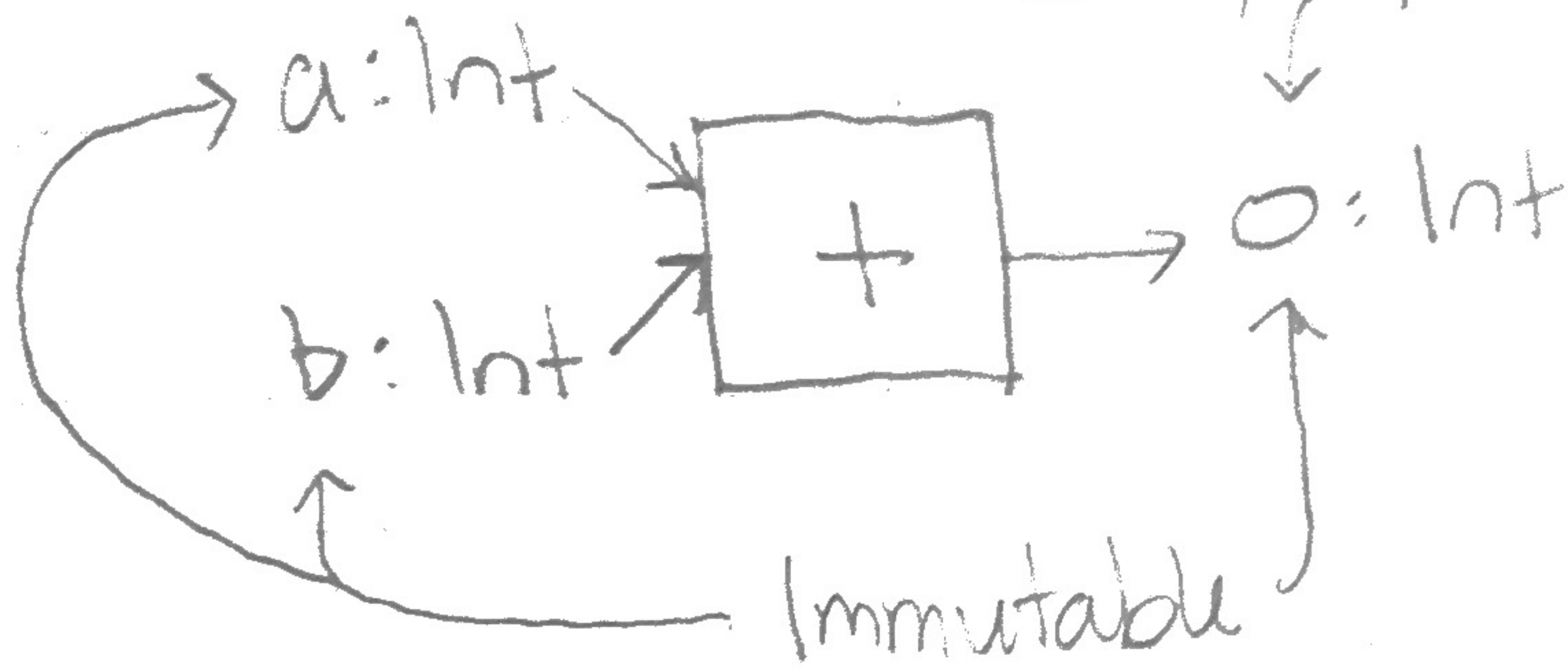


Procedure

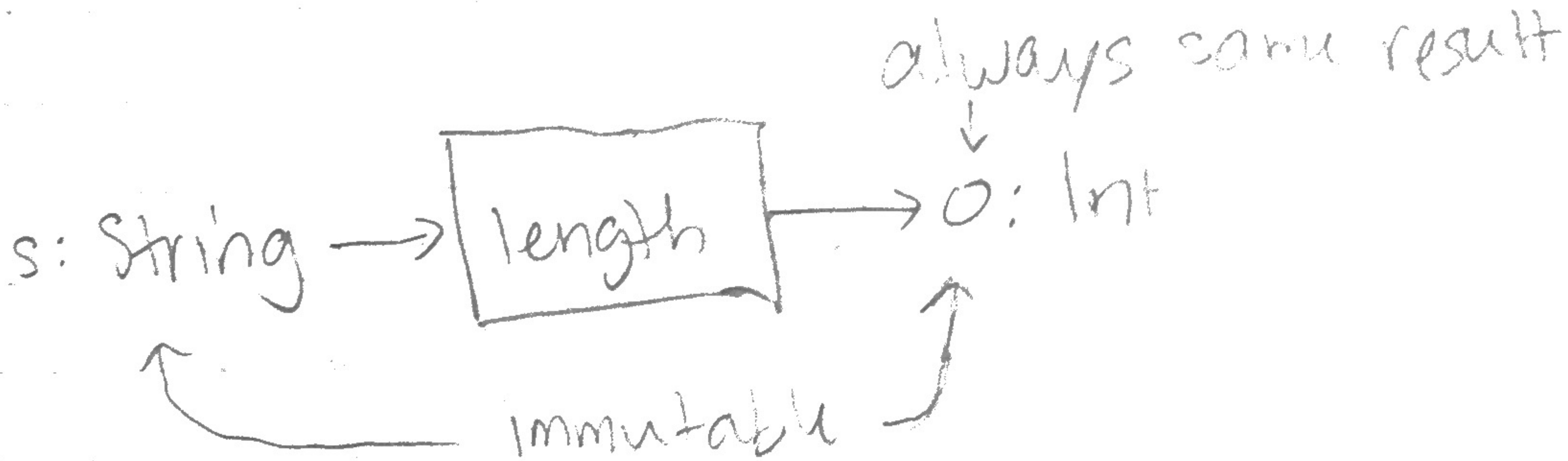


Examples (pure)

①

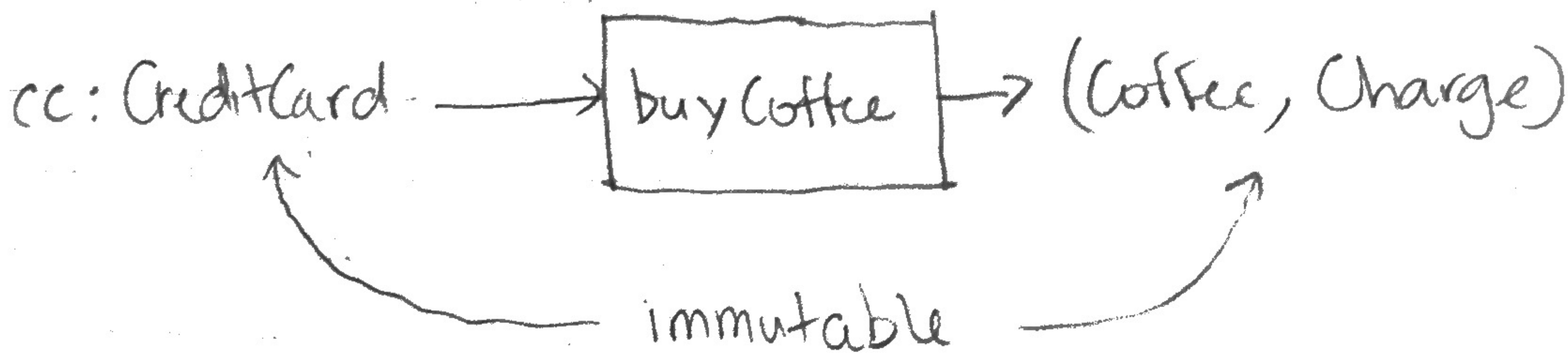


②



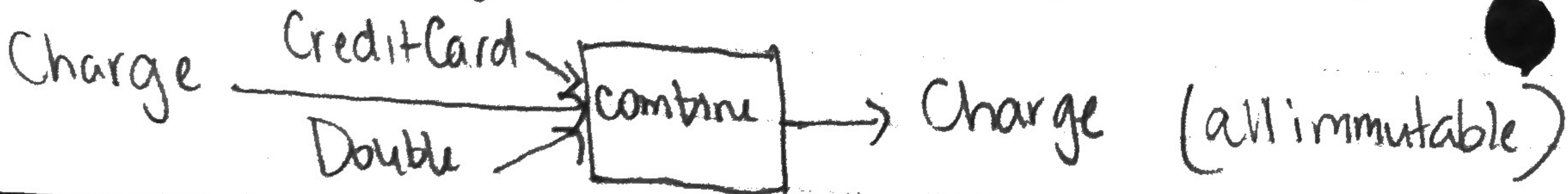
```

③ def buyCoffee(cc: CreditCard): (Coffee, Charge) = {
    val cup = new Coffee
    (cup, Charge(cc, cup.price))
}
    
```



```

④ class Charge(cc: CreditCard, amt: Double) {
    def combine(o: Charge): Charge =
        if (cc == o.cc)
            Charge(cc, amt + o.amt)
}
    
```



Because these are pure functions, they can be easily used anywhere with no impact to the rest of the system — modular + composable

↳ this is known as referential transparency (RT).

(A property of expressions in general not just functions.)

(An expression is any part of your program that can be evaluated to a result)

$\text{eval}(\text{exp}) \rightarrow \text{result}$

expression
5 > 2 + 3
↳ 5: Int ← result

The evaluation of this expression results in the same value 5 every time!

$\text{eval}(2+3)$
 $\text{eval}(\text{eval}(2) + \text{eval}(3))$
 $\text{eval}(2 + \text{eval}(3))$
 $\text{eval}(2 + 3)$
 $\text{eval}(5)$
5

This is what it means to be referentially transparent!

↳ (in fact, if we saw 2+3 in a program we could replace it with 5 and it wouldn't change the meaning of our program)

Referential Transparency + Purity

defn: An expression e is RT if, for all programs p , all occurrences of e in p can be replaced by the result of evaluating e without affecting the meaning of p .

A function f is pure if the expression $f(x)$ is RT \forall RT x .

RT \rightarrow everything a function does is represented by the value it returns.

This enables a natural mode of reasoning about program evaluation called the substitution model

(Computation proceeds like we'd solve an algebraic equation)

Example: Boolean Expressions

- (1) $! \text{true} \rightarrow \text{false}$
 - (2) $! \text{false} \rightarrow \text{true}$
 - (3) $\text{true} \ \&\& \ e \rightarrow e$
 - (4) $\text{false} \ \&\& \ e \rightarrow \text{false}$
 - (5) $\text{true} \ \|\ e \rightarrow \text{true}$
 - (6) $\text{false} \ \|\ e \rightarrow e$
- } short-circuit evaluation

$((T \ \&\& \ !T) \ \|\ (!F \ \|\ T)) \ \&\& \ ((T \ \|\ F) \ \|\ !F)$

Is this T or F?
which rules are applied?

How do we add if/else to our rules?

(7) if (true) e_1 else $e_2 \rightarrow e_1$

(8) if (false) e_1 else $e_2 \rightarrow e_2$

Can you define a rule for $\&\&$ and $\|\|$ without using $\&\&$ and $\|\|$?

$e_1 \&\& e_2 \rightarrow \text{if} (!e_1) \text{ false else } e_2$
 $e_1 \|\| e_2 \rightarrow \text{if} (e_1) \text{ true else } e_2$

How about XOR? $e_1 \wedge e_2$ is true when $e_1 \neq e_2$

$e_1 \wedge e_2 \rightarrow (e_1 \&\&!e_2) \|\| (!e_1 \&\&e_2)$

you can use $\&\&$ and $\|\|$

How does this relate to Scala code?

S> val x = "Hello, World"

S> val r₁ = x.reverse

S> val r₂ = x.reverse

} r₁ and r₂ are the same

↳ what if we substitute x with expr referenced by x?

S> val r₁ = "Hello, world".reverse

S> val r₂ = "Hello, world".reverse

} r₁ == r₂ is true

↳ x is RT

How about this?

S> val x = new StringBuilder("Hello")

S> val y = x.append(", world")

S> val r1 = y.toString } r1 and r2 are the same

S> val r2 = y.toString

* `StringBuilder.append` is a side effecting function.

S> val x = new StringBuilder("Hello")

S> val r1 = x.append(", world").toString

S> val r2 = x.append(", world").toString

↑ r1 and r2, not the same!

↓
StringBuilder.append is not pure.

point: hard to reason about side effecting code because you need to understand the larger context.

Higher-order functions

New idea: functions are values

↳ just like Int, String, List, ...

Functions can be assigned to variables, stored in data structures, and passed as arguments to functions

* When writing functional programs - it is useful to write functions that accept other functions as arguments

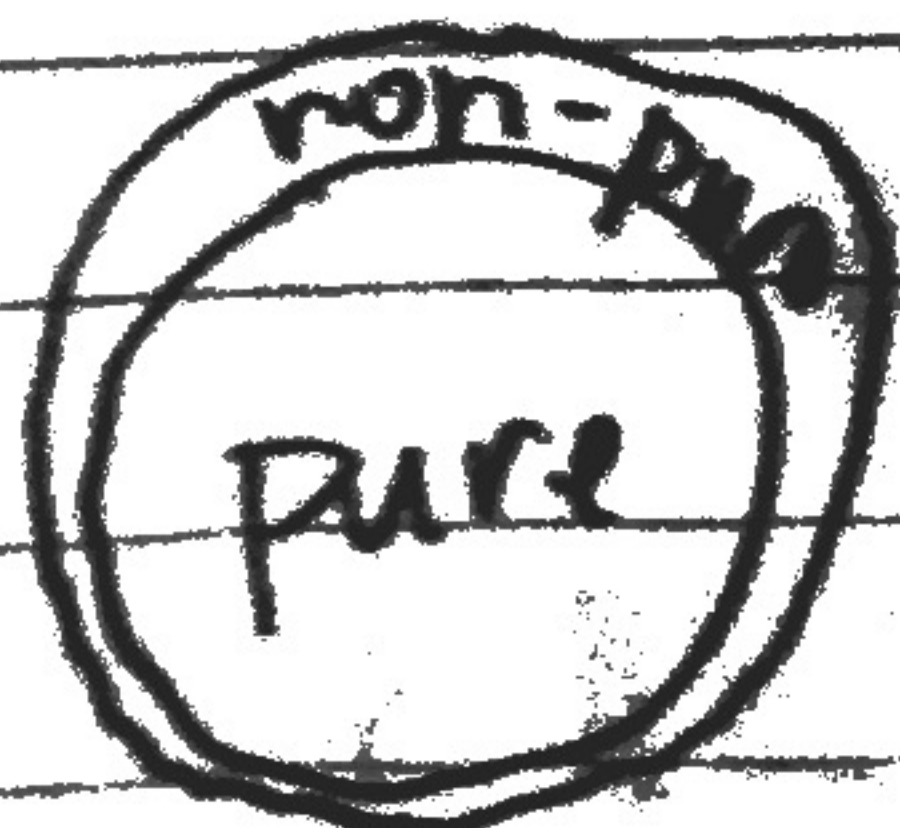
Example: higher-order / Absolute0.scala

(Now, how might we adapt this program to print both absolute value and factorial?)

Example: higher-order / Higher02.scala

Example: higher-order / Higher03.scala

Functional programs are typically written with a rich functional/pure core and a thin non-pure wrapper.



Polymorphic Functions

↳ monomorphic - functions that operate over single type.

polymorphic - functions that operate over many types.



we often recognize that we can make a function more generic (more reusable) if they can operate over many types of data.



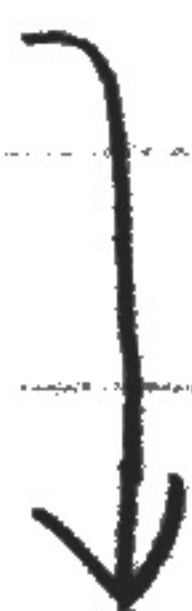
rather than writing a function for each type.

Example: `def findFirst(ss: Array[String], key: String): Int`



monomorphic

How might we "find first" over an array of type T?



polymorphic

`def findFirst[A](as: Array[A], p: A => Boolean): Int`

Example: higher-order / Poly^{OH} scala

Although, you could define function by name to pass to higher-order polymorphic functions,

it is often more more convenient to use anonymous functions.

s> (x: Int) = x == 3

↳ (Int) => Boolean

s> (x: Int, y: Int) => x + y

↳ (Int, Int) => Int

Example: higher-order / Polymos. scala

Implement the following function:

def partial1[A, B, C](a: A, f: (A, B) => C): B => C

↓

def partial1[A, B, C](a: A, f: (A, B) => C): B => C =
(b: B) => f(a, b)